

---

# Identifying Rank 2 Chvátal-Gomory Inequalities

Brady Hunsaker  
University of Pittsburgh

Craig Tovey

Ellis Johnson  
Georgia Institute of Technology  
hunsaker@engr.pitt.edu

# Outline

---

- Motivation
- Formulation of problems
- Computational issues
- User interface
- Examples
- Open-source tools used
- Future development

# Chvátal-Gomory inequalities

---

- Valid inequalities formed by using known valid inequalities.
- Technique explicitly presented by Chvátal (1973); implicit in earlier work by Gomory (1958).

1. Known valid inequalities:  $\sum_{j=1}^n a_{ij}x_j \leq b_i$ ,  $i = 1, \dots, m$ ,  
and  $x_j \geq 0$ ,  $j = 1, \dots, n$ .

2. Select nonnegative multipliers  $u_1, \dots, u_m$ .

3. The resulting Chvátal-Gomory inequality is

$$\sum_{j=1}^n \left[ \sum_{i=1}^m u_i a_{ij} \right] x_j \leq \left[ \sum_{i=1}^m u_i b_i \right].$$

# Chvátal-Gomory rank

---

Use C-G inequalities and original inequalities to generate new C-G inequalities.

- The C-G rank of the initial inequalities and any inequality dominated by a nonnegative linear combination of the initial inequalities is 0.
- Any inequality which can be generated as a C-G inequality based on rank 0 inequalities—but is not rank 0 itself—has C-G rank 1.
- A valid inequality has C-G rank  $k$  if it can be generated as a C-G inequality based on inequalities of rank  $0, 1, \dots, k - 1$ , but does not have rank less than  $k$ .

# Motivation

---

Given an IP instance and a valid inequality, what is its Chvátal-Gomory rank?

- This question is often considered when new valid inequalities are proposed for a problem.
- How much tighter is the rank 1 polyhedron than the LP relaxation?
- Determining an upper bound on the rank is often done by trial-and-error.
- Lower bound proofs are generally even harder.
- Not aware of any existing computational method for testing rank 2.

---

# Formulations and Computational Issues

# An LP for Rank 1

---

- Candidate inequality:  $\sum_{j=1}^n f_j x_j \leq g$  with  $f_j \in \mathbf{Z}, g \in \mathbf{Z}$ .
- If rank 1 C-G inequality, then there exist nonnegative  $u_i$  such that  $\sum_{i=1}^m u_i a_{ij} \geq f_j$  for all  $j$  and such that  $\sum_{i=1}^m u_i b_i < g + 1$ .
- LP for “best” multipliers:

$$\begin{array}{ll} \min & \sum_{i=1}^m b_i u_i \\ \text{s.t.} & \sum_{i=1}^m u_i a_{ij} \geq f_j \quad j = 1, \dots, n \\ & u_i \geq 0 \quad i = 1, \dots, m \end{array}$$

- If the optimal objective value satisfies  $z \leq g$ , then it is rank 0. If  $g < z < g + 1$ , then the inequality is rank 1. Otherwise it is rank  $> 1$  or invalid.

# Column generation for Rank 2

---

- Can form a similar LP for rank 2, but need a multiplier variable for each C-G inequality of rank 1.
- Use the same LP with the following column generation for a rank 1 inequality  $\sum_{j=1}^n \alpha_j x_j \leq \beta$ .  $\pi$  is the vector of dual prices for the base LP.

$$\begin{aligned} \max \quad & \sum_{j=1}^n \pi_j \alpha_j - \beta \\ \text{s.t.} \quad & \sum_{i=1}^m u_i a_{ij} = \hat{\alpha}_j \quad j = 1, \dots, n \\ & \sum_{i=1}^m u_i b_i = \hat{\beta} \\ & \hat{\alpha}_j \geq \alpha_j \quad j = 1, \dots, n \\ & [\hat{\beta}] \leq \beta \\ & u_i \geq 0 \quad i = 1, \dots, m \\ & \alpha_j, \beta \text{ integer} \end{aligned}$$

- 
- But this is not a linear IP.

# Column generation for Rank 2

---

- $[\hat{\beta}] \leq \beta$  is the same as  $\hat{\beta} - 1 < \beta$
- Replace that with  $\hat{\beta} - 1 + \epsilon \leq \beta$ , where  $\epsilon$  is small

$$\begin{aligned} \max \quad & \sum_{j=1}^n \pi_j \alpha_j - \beta \\ \text{s.t.} \quad & \sum_{i=1}^m u_i a_{ij} = \hat{\alpha}_j \quad j = 1, \dots, n \\ & \sum_{i=1}^m u_i b_i = \hat{\beta} \\ & \hat{\alpha}_j \geq \alpha_j \quad j = 1, \dots, n \\ & \hat{\beta} - 1 + \epsilon \leq \beta \\ & u_i \geq 0 \quad i = 1, \dots, m \\ & \alpha_j, \beta \text{ integer} \end{aligned}$$

# Optimization over Rank 1

---

- Think of  $\pi$  as a point in  $\mathbb{R}^n$ . Then we may view the column generation as separation.
- This allows us to separate over rank 1 inequalities.

# Certificates of rank

---

| Rank     | Certificate   |
|----------|---|
| $\leq 1$ | Multipliers for rank 0 inequalities   |
| $> 1$    | Point $x$ that satisfies rank 0 inequalities and $\sum_{j=1}^n f_j x_j \geq g + 1$  |
| $\leq 2$ | Multipliers for rank 0 and rank 1 inequalities; also certificates of rank $\leq 1$ for each rank 1 inequality used  |
|          | <ul style="list-style-type: none"><li>• Each of the certificates above is succinct.</li><li>• Unfortunately, the “certificate” for rank <math>&gt; 2</math> has to prove an IP infeasibility.</li></ul> |

# Tolerances & bounds are important

---

Potential numerical trouble due to the interaction of  $\epsilon$ , the particular instance, and the tolerances of the solver.

- Although the epsilon was written as an absolute epsilon, it is really a relative epsilon.
- May need to try several values of  $\epsilon$ .
- Need upper bounds on the multipliers  $w_i$ : no multiplier needs to be greater than 1.

---

# User Interface and Examples

# Input may be AMPL or MPS

---

Two inputs are necessary:

**IP instance** This defines the rank 0 constraints. Two possible formats:

- GNU MathProg (subset of AMPL)
- MPS

**Constraints to test** Specified in a text file containing a vector of constraint coefficients on each line.

Output is text, including certificates of rank.

# Example: High rank

---

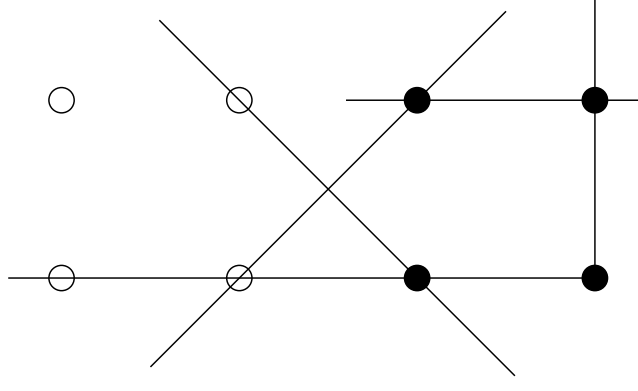
Textbook example:

$$tx_1 + x_2 \leq 1 + t$$

$$-tx_1 + x_2 \leq 1$$

$$x_1 \leq 1$$

$$x_1, x_2 \geq 0$$



Valid inequality:  $x_2 \leq 1$  has rank  $t$ .

# High rank: $t = 1$ output

```
Coin0001I At line 7 NAME          classic1
Coin0001I At line 8 ROWS
Coin0001I At line 13 COLUMNS
Coin0001I At line 17 RHS
Coin0001I At line 20 ENDATA
Coin0002I Problem classic1 has 3 rows, 2 columns and 5 elements
Coin0008I classic1 read with 0 errors
=====
Testing constraint: 0 1 <= 1
Clp0005I 0 Objective 0 Primal infeas 1 (1)
Clp0005I 2 Objective 1.5
Clp0000I Optimal - objective value 1.5
Rank of inequality is 1
Multipliers to achieve rank 1:
0.5: 1 1 2 1
0.5: 1 -1 2 1
```

# High rank: $t = 2$ output excerpts

Testing constraint:  $0 \leq 1 \leq 1$

Rank of inequality is at least 2

The following point satisfies all rank 0 inequalities and proves rank 2 or greater:

1 0.5 2 2

Adding row:  $1 \leq 1 \leq 2$

Multipliers to create rank 1 row:

0.874997000012: 1 2 2 1

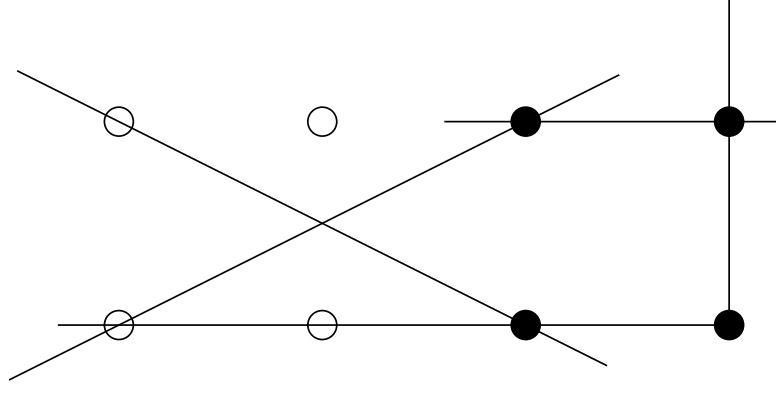
0.374997000012: 1 -2 2 1

Rank 2

Multipliers to achieve rank 2:

0.333333333333: 1 -2 2 1

0.666666666667: 1 1 2 1



# High rank: $t = 2$ output excerpts

Testing constraint:  $0 \leq 1 \leq 1$

Rank of inequality is at least 2

The following point satisfies all rank 0 inequalities and proves rank 2 or greater:

1 0.5 2 2

Adding row:  $1 \leq 1 \leq 2$

Multipliers to create rank 1 row:

0.874997000012: 1 2 2 1

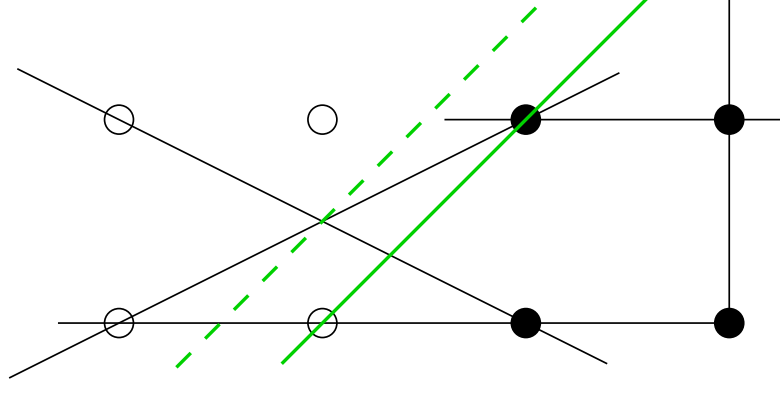
0.374997000012: 1 -2 2 1

Rank 2

Multipliers to achieve rank 2:

0.333333333333: 1 -2 2 1

0.666666666667: 1 1 2 1



# High rank: $t = 3$ output excerpts

Testing constraint:  $0 \leq 1 \leq 1$

Rank of inequality is at least 2

The following point satisfies all rank 0 inequalities and proves rank 2 or greater:

1 0.5 2 2.5

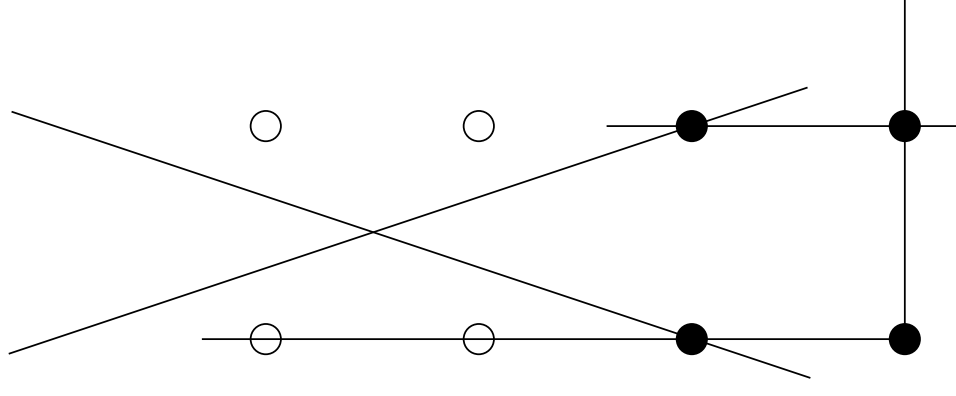
Adding row:  $0 \leq 1 \leq 2$

Multipliers to create rank 1 row:

0.59999760001: 1 3 2 1

0.59999760001: 1 -3 2 1

Rank 3 or more



# Example: Knapsack

$$\sum_{i=1}^{12} ix_i \leq 12$$
$$x_i \geq 0$$

| Facet Coefficients | RHS | C-G Rank |
|--------------------|-----|----------|
| 00001222222        | 2   |          |
| 000011122222       | 2   |          |
| 00022344666        | 6   |          |
| 001122233444       | 4   |          |
| 00123334566        | 6   |          |
| 0034567891212      | 12  |          |
| 011223445566       | 6   |          |
| 013446889111212    | 12  |          |

# Example: Knapsack

$$\sum_{i=1}^{12} ix_i \leq 12$$

$$x_i \geq 0$$

| Facet Coefficients | RHS | C-G Rank |
|--------------------|-----|----------|
| 0000122222         | 2   | 2        |
| 0000112222         | 2   | 2        |
| 0002234466         | 6   | 3+       |
| 0011222344         | 4   | 1        |
| 0012333456         | 6   | 2        |
| 00345678912        | 12  | 3+       |
| 011223445566       | 6   | 1        |
| 01344688911        | 12  | 3+       |

# Open-source tools used

---

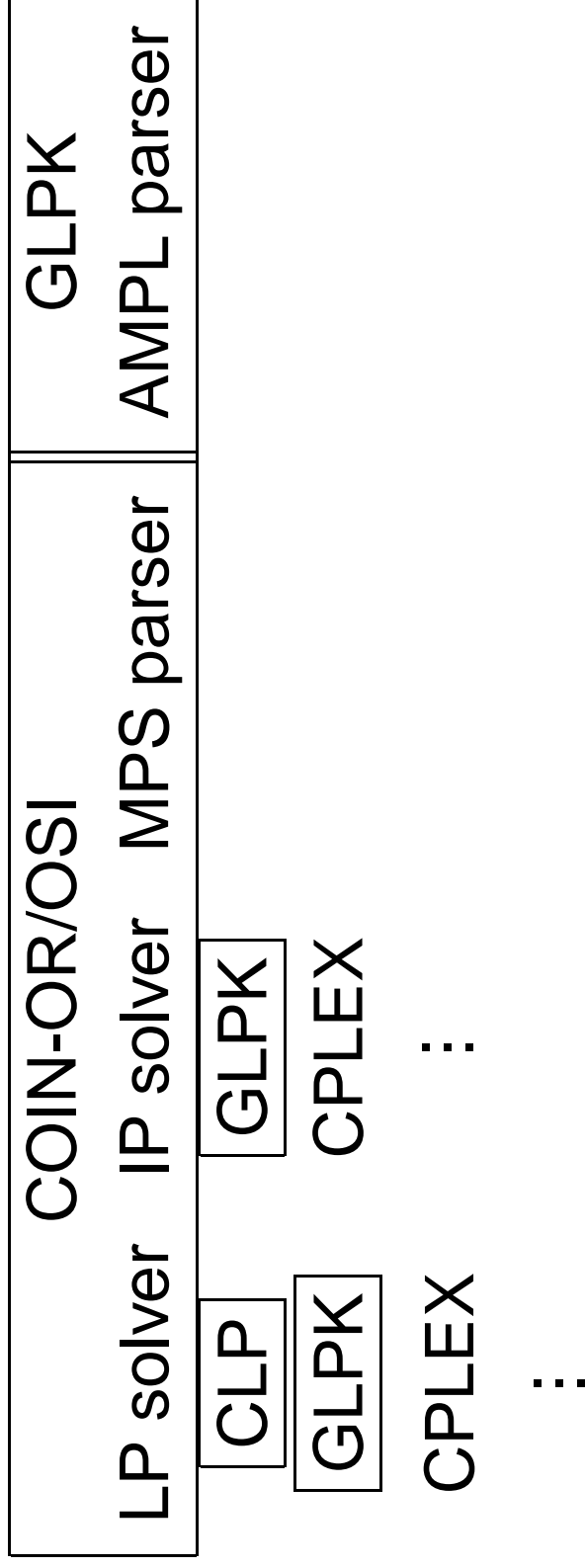
CG-rank

LP solver   IP solver   MPS parser   AMPL parser

# Open-source tools used

---

CG-rank



# Availability and license

---

- Source code will be available at <http://www.engr.pitt.edu/hunsaker/cg-rank/>
- Licensed under GNU GPL.
- Also need to download COIN-OR/OSI and GLPK in order to compile
- Email [hunsaker@engr.pitt.edu](mailto:hunsaker@engr.pitt.edu) for assistance

# Future work

---

- Improve column generation to allow for larger instances
- Allow test constraints to be specified in MPS or GNU MathProg
- Make linking with GLPK optional
- Optimization over rank 1 inequalities
- Extend to rank 3